# Strong Masting Conjecture for Multiple Size Hexagonal Tessellation in GSM Network Design 

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#### Abstract

One way to improve cellular network performance is to use efficient handover method and design pattern among other factors. The efficient design pattern has been proven geometrically to be hexagonal (Hales, 2001, pp. 1- 22) due to its maximum tessellable area coverage. But uneven geographical distribution of subscribers requires tessellable hexagons of different radii due to variation of costs of GSM masts. This will call for an overlap difference. The constraint of minimum overlap difference for multiple cell range is a new area that is untapped in cell planning. This paper addresses such multiple size hexagonal tessellation problem using a conjecture. Data from MTN River State-Nigeria, was collected. Multiple Size Hexagonal Tessellation Model (MSHTM) conjecture for masting three (3) different size MTN GSM masts in River State, accounted for least overlap difference with area of $148.3 \mathrm{~km}^{2}$ using 36 GSM masts instead of the original $21.48 \mathrm{~km}^{2}$ for 50 GSM masts. Our conjecture generally holds for k -different $(k \geq 2)$ cell range.


Keywords: disks, frequency, GSM masts, hexagon, overlap difference, tessellation.

## 1. Introduction

Cell planning is the most significant operations in GSM design network. It includes the choice of design pattern (triangular, square or hexagon), geographic, environmental and network parameters such as terrain and artificial structures, base station location and transmission power among others. But the hexagonal design has least overlap and hence has a strength higher than both a square and an equi-triangular polygon. The hexagon motivated circular shaped cells are produced when two or more sector signal radiated antenna are used. These circular cells overlap significantly and is crucial for subscriber movement-handover. This overlap removes signal loss due to no coverage or ensures soft handover due to better overlapping regions. How much overlap difference to be permitted in the uniform design network has been studied extensively by Donkoh E. K. et.al (2015a). Uneven geographical distribution of subscribers require varied cell ranges. This will result in variation in overlap difference. Both the design pattern and geometry of overlap is complex but the result is more economical as fewer GSM masts will be used. Realistically, how much overlap has been a thriving challenge in recent times due to variation in subscriber geographical distribution, affordability and significance of mobile telephone in this age of technology.

## 2. Related Literature

Antenna signal radiation in GSM network design takes many form. But the most profitable radiation pattern is known to be the sector motivated circular shape (Azad, 2012.). Due to it convenience and usefulness manufacturers of GSM antenna like Mobile Mark, Multiband Technologies, Global Source, Asian Creation among others have designed GSM antenna's with variety of sector angles including $30^{\circ}, 60^{\circ}, 90^{\circ}, 120^{\circ}$ as the common ones. Nonetheless, designing GSM network due to uneven geographical distribution of subscribers requires geometry of 2-D hexagonal tessellation . Unfortunately, multiple size hexagonal design has since not been generally practiced due to the complexity in the position of the base station that offers minimum overlap difference. This paper uses the least tessellable polygonal overlap difference to conjecture a multiple size design formula for hexagons of several dimensions.
Donkoh \& Opoku (2016, pp.33) emphasize with geometric proof that the hexagon has the least overlap difference of $13.4 \%$ and hence has the strongest tessellable area coverage.

Donkoh E. K et.al (2015a, pp. 5) investigated the hexagonal overlap difference for a uniform 0.6 km cell range of 50 GSM masts design and obtain an overlap difference of 5.788 km using 35 GSM masts instead of the original 26.884 km .

This however confirms the economies of scale for design pattern minimizing overlap difference with smooth handover.
Hamad-Ameen J.J (2008, pp. 393), studied frequency and cell planning with clusters ( $\mathrm{K}=7,9,13$ ) and without clusters for uniform cell range and concluded that it is very efficient in GSM masting.
Donkoh E. K et. al (2015c), the authors considered GSM network design with two different cell range namely 1 km and 3 km obtaining the least overlap difference of 15.27 km as compared to the original 30.95 km .
In application to earth and mineral science, Raposo (2011), uses vertex-clustering of uniform hexagonal tessellation to simplify scale-specific automated map line.
Beyond applications of Christaller (1933) classic theory, hexagonal tessellation has been advocated for thematic cartography by Carr et.al (1992), and has been used to study cluster perception in animated maps (Griffin, 2006), as well as color perception (Brewer, 1996).
However, geometry of multiple size hexagonal tessellation in covering bounded areas, optimizing the overlap difference for k-different ( $k \geq 2$ ) cell ranges have not been studied. We therefore proposed the strong multiple size masting conjecture for solving the bounded area coverage problem in GSM network design.

## 3. Computational Experience

We consider intersecting circular cells superimposed on non-hexagonal polygon with cell range $R_{1}$ and $R_{2}$ and ${ }^{1}$ apothems $\mathrm{r}_{1}$ and $r_{2}$ respectively. In Figure 1, XZ is one side of the non-hexagonal polygon and the cells overlap for covering purposes with an overlap difference(d) of $R_{2}-r_{2}+R_{1}-r_{1}$.


Figure 1. Overlap difference for Non-Uniform cell Range $R_{1}, R_{2}$.
From triangle $X O Y$ in Figure 1, $R_{2}^{2}=h^{2}+r_{2}^{2}$ where $X Y$ is a side of the right triangle.

$$
\begin{align*}
h & =\sqrt{R_{2}^{2}-r_{2}^{2}}  \tag{1}\\
\sin \theta_{2} & =\frac{h}{R_{2}} \\
h & =R_{2} \sin \theta_{2} \tag{2}
\end{align*}
$$

Equating (1) and (2),

$$
\begin{align*}
R_{2} \sin \theta_{2} & =\sqrt{R_{2}^{2}-r_{2}^{2}} \\
\sin \theta_{2} & =\sqrt{1-\left(\frac{r_{2}^{2}}{R_{2}^{2}}\right)} \tag{3}
\end{align*}
$$

Also, $r_{2}=R_{2} \cos \theta_{2}$
Similarly in triangle $X O^{\prime} Y, R_{1}^{2}=h^{2}+r_{1}^{2}$ where $h$ is a side of the right triangle.

$$
\begin{align*}
h & =\sqrt{R_{1}^{2}-r_{1}^{2}}  \tag{4}\\
\sin \theta_{1} & =\frac{h}{R_{1}} \\
h & =R_{1} \sin \theta_{1} \tag{5}
\end{align*}
$$

Equating (4) and (5),

[^0]\[

$$
\begin{align*}
R_{1} \sin \theta_{1} & =\sqrt{R_{1}^{2}-r_{1}^{2}} \\
\sin \theta_{1} & =\sqrt{1-\left(\frac{r_{1}^{2}}{R_{1}^{2}}\right)} \tag{6}
\end{align*}
$$
\]

Similarly in triangle $\mathrm{XO}^{\prime} \mathrm{Y}, \mathrm{R}_{1}^{2}=\mathrm{h}^{2}+\mathrm{r}_{1}^{2}$ where h is a side of the right triangle.
Also, $r_{1}=R_{1} \cos \theta_{1}$
Equating (2) and (5) and extending it to radius $R_{2}$

$$
\begin{align*}
R_{1} \sin \theta_{1} & =R_{2} \sin \theta_{2} \\
R_{1} & =R_{2} \frac{\sin \theta_{2}}{\sin \theta_{1}} \\
R_{1} & =\frac{R_{1} R_{2} \sin \theta_{2}}{\sqrt{R_{1}^{2}-r_{1}^{2}}} \\
\sin \theta_{2} & =\frac{\sqrt{R_{1}^{2}-r_{1}^{2}}}{R_{2}} \tag{7}
\end{align*}
$$

Similarly, in triangle $X O^{\prime} Y$,

$$
\begin{equation*}
\sin \theta_{1}=\frac{\sqrt{R_{2}^{2}-r_{2}^{2}}}{R_{1}} \tag{8}
\end{equation*}
$$

A single overlap difference for Figure 2 is

$$
\begin{align*}
& d_{1}=R_{2}-r_{2}+R_{1}-r_{1}  \tag{9}\\
& d_{1}=\sum_{i=1}^{2} R_{i}-r_{i} \\
& d_{1}=\sum_{i=1}^{2} R_{i}\left(1-\cos \theta_{i}\right) \tag{10}
\end{align*}
$$

Generally n non-uniform overlaps,

$$
\begin{align*}
& d_{n}=\sum_{i=1}^{2 n} R_{i}-r_{i} \\
& d_{n}=\sum_{i=1}^{2 n} R_{i}\left(1-\cos \theta_{i}\right) \tag{11}
\end{align*}
$$

Equation (9), (10) or (11) is used to calculate the overlap difference for non-uniform disks as shown in Figure 1. We established a formula for calculating the area of a pair of overlap for non-hexagonal polygon inscribed disks. From Figure 1 we have:
Area of single overlap difference $=$ (Area of sector $X O^{\prime} Z-$ Area of triangle $X O^{\prime} Z$ ) + (Area of sector $X O Z-$ Area of triangle $X O Z$ )

$$
\begin{equation*}
A_{d}=\frac{1}{2} R_{1}^{2}\left(\theta_{1}-\sin \theta_{1}\right)+\frac{1}{2} R_{2}^{2}\left(\theta_{2}-\sin \theta_{2}\right) \tag{12}
\end{equation*}
$$

Substituting (7) and (8) into (12)

$$
\begin{equation*}
A_{d}=\frac{1}{2} R_{1}^{2}\left[\sin ^{-1}\left(\frac{\sqrt{R_{2}^{2}-r_{2}^{2}}}{R_{1}}\right)-\frac{\sqrt{R_{2}^{2}-r_{2}^{2}}}{R_{1}}\right]+\frac{1}{2} R_{2}^{2}\left[\sin ^{-1}\left(\frac{\sqrt{R_{1}^{2}-r_{1}^{2}}}{R_{2}}\right)-\frac{\sqrt{R_{1}^{2}-r_{1}^{2}}}{R_{2}}\right] \tag{13}
\end{equation*}
$$

Generally for n different overlaps,

$$
\begin{align*}
& A_{n}=\frac{1}{2} \sum_{i=1}^{2 n} R_{i}^{2}\left(\theta_{i}-\sin \theta_{i}\right) \\
& A_{d}=\frac{1}{2} \sum_{i=1}^{2 n} R_{i}^{2}\left[\sin ^{-1}\left(\frac{\sqrt{R_{i+1}^{2}-r_{i+1}^{2}}}{R_{i}}\right)-\frac{\sqrt{R_{i+1}^{2}-r_{I+1}^{2}}}{R_{i}}\right] \tag{14}
\end{align*}
$$

Equation (13) is the area of each non-uniform cell range in terms of overlap difference that is not created by hexagon. The value $A_{d}$ can be calculated from Figure 1 in Autocad environment or using equation (13) as shown in column 4 and 8 in Table 2 where $R_{1}$ and $R_{2}$ are the cell ranges of the GSM masts. The equation $d=d_{m}-d_{n}$ is the overlap difference between disks with centres m and n .

### 3.1 Overlap for Optimal Disks Covering in Hexagonal Tessellation.

Donkoh E. K et.al (2015b, pp.26) in their study of overlap dimensions in cyclic tessellable regular polygon emphasize that hexagon has better overlap difference of $13.7 \%$ as compared to $29.3 \%$ for square and $50.0 \%$ for equi-triangular polygon.

$$
\text { Overlap difference }= \begin{cases}\text { Hexagon, } \quad 0<\text { Strong } \leq 13.7 \% \\ \text { Square }, & 29.3 \% \leq \text { Moderate }<50 \% \\ \text { Equi }- \text { triangular }, & \text { Weak } \geq 50 \%\end{cases}
$$

Figure 2(a) illustrates the hexagonal cell layout with inradius and circumradius of the hexagonal cell as $r_{1}$ and $R_{1}$, respectively. In Figure 2(b), cells partially overlapped because $R_{1}$ equals to the hexagon's circumradius.

(a) Hexagonal cell layout

(b) Idealized circular layout

Figure 2. Cell Layout Model's For GSM Networks.

### 3.2 Generalized Overlap Difference in Hexagon-Inscribed Disks.

Overlap in cell planning ensures smooth handover in GSM network with a differential effect such as increasing the number of GSM masts, environmental destruction etc and therefore must be kept as minimal as possible. A typical overlap may arise as a result of uniform or non-uniform cell range. Ample research work has been done on the uniform cell range (Donkoh E. K, et al 2015a). We minimize the overlap difference for non-uniform cell range of two different GSM antenna masts with radii $R_{1}$ and $R_{2}\left(R_{2}>R_{1}\right)$ and corresponding apothem $r_{1}$ and $r_{2}$ to be $R_{1}-r_{1}+2 R_{2}-r_{2}$. A mixture of non-uniform cell radius with wrong frequency assignment results in large overlap difference and in effect increases interference (such as cross talk, background noise, error in digital signaling-missed calls, blocked calls, dropped calls, etc). It also has the differential advantage of reducing cost of coverage area as multiple cell range are factored in the design. Figure 4 shows 1 : k side length, for $\mathrm{k} \in \mathbb{N} \geq 2$ overlap difference for two different radii $R_{1}$ and $R_{2}$.


Figure 4. Overlap difference for non-uniform disks, represented by ${ }^{\prime \prime}{ }^{-",}$, and GSM cell by O .

Let $d_{R_{2}, R_{1}}=d+d_{1}+d_{2}$ represents the global least overlap difference for two different size disks superimposed on tessellable hexagon in the ratio $1: \mathrm{k}$, where

$$
\begin{align*}
& d_{R_{2}, R_{1}}=R_{1}-r_{1}+R_{1}+R_{2}-r_{2} \\
& d_{R_{2}, R_{1}}=R_{2}-r_{2}+2 R_{1}-r_{1} \tag{18}
\end{align*}
$$

Theorem: The apothem $r_{n}$ created by $n$ sided tessellable regular polygon inscribed in a disk of radius $R_{1}$ is

$$
r_{n}=R_{1} \operatorname{Cos}\left(\frac{\pi}{n}\right)
$$

Proof.
Donkoh et al (2016, pp.33-34) gave a formal proof of this theorem. Equation (18) then becomes

$$
\begin{align*}
& d_{R_{2}, R_{1}}=R_{2}-R_{2} \cos \left(\frac{\pi}{n}\right)+2 R_{1}-R_{1} \cos \left(\frac{\pi}{n}\right) \\
& d_{R_{2}, R_{1}}=R_{2}\left[1-\cos \left(\frac{\pi}{n}\right)\right]+R_{1}\left[2-\cos \left(\frac{\pi}{n}\right)\right] \tag{19}
\end{align*}
$$

Since the polygon is a hexagon $n=6$ sided but with different radii. Thus

$$
\begin{align*}
& d_{R_{2}, R_{1}}=R_{2}\left[1-\cos \left(\frac{\pi}{6}\right)\right]+R_{1}\left[2-\cos \left(\frac{\pi}{6}\right)\right] \\
& d_{R_{2}, R_{1}}=\frac{1}{2}\left[(2-\sqrt{3}) R_{2}+(4-\sqrt{3}) R_{1}\right] \tag{20}
\end{align*}
$$

Equation (20) is the least overlap difference for GSM network design using two different radii since the 1 : $n$ size hexagon tile completely.

### 3.3 Masting Conjecture

Generally for $k$ different tessellable regular polygons $n_{1}, n_{2}, \ldots n_{k}$ inscribed in disks with respective radii $R_{k}, R_{k-1}, \ldots, R_{1}$ (where $R_{k}>R_{k-1}$ ) the least overlap difference is

$$
\begin{align*}
& d_{n_{k}, n_{k-1}, \ldots n_{1}}=R_{k}\left[1-\cos \left(\frac{\pi}{n_{1}}\right)\right]+R_{k-1}\left[2-\cos \left(\frac{\pi}{n_{2}}\right)\right]+\cdots+R_{1}\left[k-\cos \left(\frac{\pi}{n_{k}}\right)\right] \\
& d_{n_{k}, n_{k-1}, \ldots n_{1}}=\sum_{l=1}^{k} R_{k-l+1}\left[l-\cos \left(\frac{\pi}{n_{k}}\right)\right] \tag{21}
\end{align*}
$$

Since we are tilling with hexagon $n_{k}=6, \forall k \in \mathbb{R}$

$$
\begin{align*}
& d_{n_{k}, n_{k-1}, \ldots n_{1}}=\sum_{l=1}^{k} R_{k-l+1}\left[l-\cos \left(\frac{\pi}{6}\right)\right] \\
& d_{n_{k}, n_{k-1}, \ldots n_{1}}=\sum_{l=1}^{k} R_{k-l+1}\left[l-\frac{\sqrt{3}}{2}\right] \tag{22}
\end{align*}
$$

For three different tessellable regular hexagon the least overlap difference can be obtain from equation (22) to be

$$
\begin{align*}
d_{R_{3}, R_{2}, R_{1}} & =\sum_{l=1}^{k=3} R_{k-l+1}\left[l-\frac{\sqrt{3}}{2}\right] \\
& =R_{3}\left(1-\frac{\sqrt{3}}{2}\right)+R_{2}\left(2-\frac{\sqrt{3}}{2}\right)+R_{1}\left(3-\frac{\sqrt{3}}{2}\right) \\
& =R_{3}-R_{3} \frac{\sqrt{3}}{2}+2 R_{2}-R_{2} \frac{\sqrt{3}}{2}+3 R_{1}-R_{1} \frac{\sqrt{3}}{2} \\
& =R_{3}-R_{3} \cos \left(\frac{\pi}{6}\right)+2 R_{2}-R_{2} \cos \left(\frac{\pi}{6}\right)+3 R_{1}-R_{1} \cos \left(\frac{\pi}{6}\right) \\
d_{n_{3}, n_{2}, n_{1}} & =R_{3}-r_{3}+2 R_{2}-r_{2}+3 R_{1}-r_{1} \tag{23}
\end{align*}
$$

This is the one sided least overlap difference of triple non-uniform hexagonal tessellation for masting in GSM network. Equation (23) significantly informs cell planners how to design the GSM network for least overlap difference. Figure 6 illustrate this concept.


Figure 6. Triple Size Hexagonal Tessellation
Continuous upward tiling will lead us to the following overlap differences .
Case I: Hexagon with Radius $R_{2}$ as in 6(b)
Overlap difference will be

$$
\begin{aligned}
& d=2\left(R_{2}-r_{2}\right) \\
& d=(2-\sqrt{3}) R_{2}
\end{aligned}
$$

For $k$ overlaps, the difference will be

$$
\begin{equation*}
d_{k}=(2-\sqrt{3}) k R_{2} \tag{24}
\end{equation*}
$$

Case II: Hexagons with radii $R_{2}$ and $R_{1}$


6(c): Section of hexagon tilling with radii $R_{2}, R_{1}$


6(d): Section of hexagon tiling with radii $R_{2}, R_{1}$ Overlap difference will be

$$
\begin{aligned}
& d=R_{2}-r_{2}+R_{1}-r_{1} \\
& d=R_{2}\left(1-\cos \frac{\pi}{6}\right)+R_{1}\left(1-\cos \frac{\pi}{6}\right) \\
& d=\frac{1}{2}(2-\sqrt{3})\left(R_{1}+R_{2}\right)
\end{aligned}
$$

For $m$ such overlaps, the difference will be

$$
\begin{equation*}
d_{m}=\frac{1}{2}(2-\sqrt{3})\left(R_{1}+R_{2}\right) m \tag{25}
\end{equation*}
$$

Case III: Hexagons with Radii $R_{2}$ and $R_{1}$
Overlap difference is

$$
\begin{aligned}
& d=R_{2}-r_{2}+R_{1}-r_{1} \\
& d=R_{2}\left[1-\cos \left(\frac{\pi}{6}\right)\right]+R_{1}\left[1-\cos \left(\frac{\pi}{6}\right)\right] \\
& d=\frac{1}{2}(2-\sqrt{3})\left(R_{1}+R_{2}\right)
\end{aligned}
$$

For $p$ such overlaps, the difference will be

$$
\begin{equation*}
d_{p}=\frac{1}{2}(2-\sqrt{3})\left(R_{1}+R_{2}\right) p \tag{26}
\end{equation*}
$$

Case IV: Hexagons with Radius $R_{1}$


Figure 6(e). Section of hexagonal tilling with radius $R_{1}$
Overlap difference will be

$$
\begin{aligned}
& d=2\left(R_{1}-r_{1}\right) \\
& d=2\left(R_{1}-R_{1} \cos \frac{\pi}{6}\right)
\end{aligned}
$$

For $n$ such overlaps, the difference will be

$$
\begin{equation*}
d_{n}=(2-\sqrt{3}) n R_{1} \tag{27}
\end{equation*}
$$

Generally, the various cases put together gives us total overlap difference

$$
\begin{align*}
& d=1\left(R_{3}-r_{3}+2 R_{2}-r_{2}+3 R_{1}-r_{1}\right)+(2-\sqrt{3}) k R_{2}+\frac{1}{2}(2-\sqrt{3})\left(R_{1}+R_{2}\right) p+\frac{1}{2}(2-\sqrt{3})\left(R_{1}+R_{2}\right) m+(2-\sqrt{3}) n R_{1} \\
& d=\left[R_{3}\left(1-\frac{\sqrt{3}}{2}\right)+R_{2}\left(2-\frac{\sqrt{3}}{2}\right)+R_{1}\left(3-\frac{\sqrt{3}}{2}\right)\right]+(2-\sqrt{3})\left(k R_{2}+n R_{1}\right)+\frac{1}{2}(2-\sqrt{3})\left(R_{1}+R_{2}\right)(m+p) \tag{28}
\end{align*}
$$

Table 1. WGS-84 coordinates of 50 MTN GSM masts in River State, Nigeria.



| 49. | CHIEF AKAROLO ESTATE, ELEKAHIA |  | $4^{0} 49^{\prime} 04.98441^{\prime \prime}$ | N | $7^{0} 01^{\prime}$ | $29.27181^{\prime \prime}$ | E | 280937.839 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 50. | DANGOTE PREM., TRANS AMADI, OGINIGBA | $4^{0} 49^{\prime} 16.44446^{\prime \prime} \mathrm{N}$ | $7^{0} 02^{\prime} 15.95166^{\prime \prime}$ | E | 282377.492 | 533217.187 |  |  |

A plot of the 50 MTN GSM masts coordinates is shown in Figure 7(a). We write a matlab code for the covering as shown in Figure 7(b).


Figure 7(a). Position of 50 MTN GSM cell coordinates with cell band in River State, Nigeria

A Matlab code for 50 MTN GSM Masts position • has been covered with 36 hexagons $\square$, in River State-Nigeria. This is written below.

```
>>load set3
>>scatter(set3(:,1),set3(:,2),'fill')
>>n=5.1;
>>m=0.96;
>>figure(1),hold on
>>for i=5.0008:.015:n
    >>for j = 0.81325:0.015:m
        >>hexagon(0.005,i,j)
    >> end
>>end
    >>for i=5.0008:.015:n
        >>for j = 0.81325:0.015:m
    >>hexagon(0.005,i+0.0075,j+0.0075)
    >>end
    >>end
>> axis([[5 5.1 0.8 0.96])
>>xlabel('Easterns (xkm)')
>>ylabel('Northerns (ykm)')
```



Figure 7(b). Graph of GSM masts nodes
>>title('Hexagonal Tessellation of MTN RIVER STATE, NIGERIA Masts').
This is shown in Figure 7(b). We plot the Figure 7(b) on Autocad and compute the overlap difference and overlap area shown in Table 2. Figure 8 shows the plot.


Figure 8. Optimal disks covering for triple size hexagons using MTN masts, River State - Nigeria.
Table 2 shows the overlap difference ( $d$ ) and overlap area $\left(A_{d}\right)$ obtain as a result of Figure 7(b).

Table 2. Overlap difference for $0.6 \mathrm{~km}, 1.3 \mathrm{~km}$, 2.5 km MTN cell range - River State, Nigeria

| Serial | Overlap <br> Difference <br> $d=d_{m}-d_{n}$ | Value (m) | Area of <br> overlap $\left(A_{d}\right)$ | Serial | Overlap Differen <br> $d=d_{m}-d_{n}$ | Value (m) | $\left.\begin{array}{c}\text { Area of } \\ \text { overlap }\end{array} A_{d}\right)$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |


| 13. | $d_{6}-d_{5}$ | 131.4954 | 43632.25393 | 34. | $d_{44}-d_{37}$ | 453.1748 | 518224.6067 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 14. | $d_{5}-d_{14}$ | 27.5591 | 1916.534266 | 35. | $d_{44}-d_{26}$ | 206.3699 | 107468.0168 |
| 15. | $d_{40}-d_{38}$ | 131.1031 | 43372.29959 | 36. | $d_{32}-d_{1}$ | 129.4741 | 42301.16504 |
| 16. | $d_{40}-d_{17}$ | 381.1399 | 366568.6824 | 37. | $d_{32}-d_{37}$ | 292.9307 | 216529.1175 |
| 17. | $d_{43}-d_{35}$ | 633.2396 | 1011865.197 | 38. | $d_{27}-d_{1}$ | 273.7579 | 189112.3313 |
| 18. | $d_{43}-d_{13}$ | 633.2396 | 1011865.197 | 39. | $d_{34}-d_{9}$ | 487.4520 | 599584.2833 |
| 19. | $d_{9}-d_{30}$ | 418.2921 | 441515.3955 | 40. | $d_{24}-d_{25}$ | 404.3435 | 412560.3237 |
| 20. | $d_{9}-d_{48}$ | 669.0557 | 1129564.61 | 41. | $d_{24}-d_{48}$ | 483.9451 | 590988.0793 |
| 21. | $d_{9}-d_{1}$ | 464.3242 | 544037.9123 | 42. | $d_{29}-d_{21}$ | 399.7777 | 403295.761 |
| 43. | $d_{45}-d_{26}$ | 182.0447 | 83626.24686 | 63. | $d_{29}-d_{19}$ | 23.2047 | 1358.7465 |
| 44. | $d_{45}-d_{21}$ | 264.4319 | 176446.9753 | 64. | $d_{19}-d_{20}$ | 294.0752 | 218224.4106 |
| 45. | $d_{45}-d_{29}$ | 212.4948 | 113941.817 | 65. | $d_{19}-d_{21}$ | 534.0102 | 719590.8493 |
| 46. | $d_{26}-d_{37}$ | 341.6531 | 294548.8004 | 66. | $d_{19}-d_{26}$ | 13.6140 | 467.6899 |
| 47. | $d_{26}-d_{3}$ | 636.5535 | 1022483.593 | 67. | $d_{20}-d_{2}$ | 280.9723 | 199211.103 |
| 48. | $d_{26}-d_{19}$ | 13.6140 | 467.6899306 | 68. | $d_{7}-d_{33}$ | 394.9396 | 393593.4759 |
| 49. | $d_{3}-d_{21}$ | 783.8926 | 1550599.56 | 69. | $d_{33}-d_{11}$ | 624.8460 | 985218.4021 |
| 50. | $d_{3}-d_{20}$ | 99.3801 | 24922.14313 | 70. | $d_{33}-d_{39}$ | 40.3124 | 4100.7551 |
| 51. | $d_{3}-d_{19}$ | 576.9297 | 839909.1656 | 71. | $d_{39}-d_{11}$ | 448.9873 | 508691.6872 |
| 52. | $d_{3}-d_{37}$ | 100.0536 | 25261.08301 | 72. | $d_{41}-d_{38}$ | 597.6027 | 901180.1747 |
| 53. | $d_{3}-d_{45}$ | 5.3750 | 72.90267503 | 73. | $d_{41}-d_{17}$ | 305.7118 | 235836.4513 |
| 54. | $d_{25}-d_{1}$ | 298.0264 | 224127.9406 | 74. | $d_{38}-d_{40}$ | 131.1031 | 43372.2996 |
| 55. | $d_{25}-d_{48}$ | 686.8085 | 1190303.882 | 75. | $d_{18}-d_{44}$ | 517.7495 | 676434.7394 |
| 56. | $d_{25}-d_{30}$ | 288.3549 | 209817.2539 | 76. | $d_{18}-d_{23}$ | 644.1559 | 1047052.613 |
| 57. | $d_{25}-d_{20}$ | 152.7068 | 58844.14814 | 77. | $d_{18}-d_{37}$ | 345.0220 | 300386.2877 |
| 58. | $d_{24}-d_{30}$ | 712.9229 | 1282542.18 | 78. | $d_{18}-d_{26}$ | 810.5369 | 1657799.9 |
| 59. | $d_{24}-d_{7}$ | 611.8305 | 944601.809 | 79. | $d_{18}-d_{3}$ | 247.7574 | 154895.8552 |
| 60. | $d_{18}-d_{45}$ | 105.7063 | 28196.04989 | 80. | $d_{18}-d_{21}$ | 77.0424 | 14977.7347 |
| Total $\left(\sum d\right)=26,412.518 \mathrm{~m}$ |  |  |  | Total $\left(\sum A_{d}\right)=32,834,104.29 \mathrm{~m}^{2}$ |  |  |  |

Let $R_{i}$ be disks with radius,$N_{R_{i}}$ be number of cells with radius $i$ then
Area of cells $=$ sum of area of all disks-sum of all overlap areas of cells
Area $=\sum_{i=1}^{n} N_{R_{i}} \times \pi R_{i}^{2}-\sum A_{d}$
Where the overlap area is calculated using the inclusive exclusive formula

$$
\left|\bigcup_{i=1}^{n} A_{i}\right|=\text { Area } \sum_{i}\left|A_{i}\right|-\text { Area } \sum_{i<j}\left|A_{i} \cap A_{j}\right|+\text { Area } \sum_{i<j<k}\left|A_{i} \cap A_{j} \cap A_{k}\right|-\cdots+(-1)^{n+1} \text { Area }\left|\bigcap_{i=1}^{n} A_{i}\right|
$$

where the first sum is over all $i$ the second sum is over all pairs $i . j$ with $i<j$, the third sum is over all triples $i, j, k$ with $i<j<k$ and so fourth.

$$
=\left(48 \times \pi \times 600^{2}+1 \times \pi \times 1300^{2}+1 \times \pi \times 6600^{2}-32,834,104.29\right) m^{2}
$$

$$
\text { Area }=21.48 \mathrm{~km}^{2}
$$

We compute the total coverage area of the proposed multiple size hexagonal tessellation employed in the design and compare with the existing coverage area.

Total area of hexagon $=$ sum of Number of hexagons $\times$ unit area

$$
\begin{aligned}
& =\sum_{i=1}^{n} N_{R_{i}} \times \frac{3 \sqrt{3}}{2} R_{i}^{2} \\
& =\left(34 \times \frac{3 \sqrt{3}}{2} \times 600^{2}+1 \times \frac{3 \sqrt{3}}{2} \times 1200^{2}+1 \times \frac{3 \sqrt{3}}{2} \times 6600^{2}\right) \mathrm{m}^{2} \\
\text { Area } & =148.71 \mathrm{~km}^{2}
\end{aligned}
$$

Table 3. Comparative Analysis of the MSHT model to the Original Layout method for MTN River State, Nigeria.

| Case Study | Original Layout | Multiple Size Hexagonal Tessellation |
| :---: | :---: | :---: |
|  | Number ofmasts $\quad$Area covered <br> $\left(\mathrm{km}^{2}\right)$ | Numberof masts $\quad$Area covered <br> $\left(\mathrm{km}^{2}\right)$ |
| MTN Nigeria, River State $R_{1}=0.6 \mathrm{~km}, R_{2}=1.3 \mathrm{~km}, R_{3}=2.5 \mathrm{~km}$ | $50 \quad 21.48$ | $36 \quad 148.71$ |
| Number of overlaps | 80 | $\mathrm{k}=1, \mathrm{~m}=0, \mathrm{p}=1, \mathrm{n}=52$ |
| $\begin{aligned} & d=\left[R_{3}\left(1-\frac{\sqrt{3}}{2}\right)+R_{2}\left(2-\frac{\sqrt{3}}{2}\right)+R_{1}\left(3-\frac{\sqrt{3}}{2}\right)\right] \\ & +(2-\sqrt{3}) k R_{2}+\frac{1}{2}(2-\sqrt{3})\left(\mathrm{R}_{1}+\mathrm{R}_{2}\right) \mathrm{p} \\ & +\frac{1}{2}(2-\sqrt{3})\left(\mathrm{R}_{1}+\mathrm{R}_{2}\right) \mathrm{m}+(2-\sqrt{3}) \mathrm{nR}_{1} \end{aligned}$ | 26.413 km | 12.368 km |
| Ratio |  | 1mast $\quad 4.13 \mathrm{~km}{ }^{2}$ |
| Number of GSM masts based on cell range | $\begin{gathered} N_{R_{1}=0.6 \mathrm{~km}^{-}-48}, \\ N_{R_{2}=1.3 \mathrm{~km}^{-1}}, \\ N_{R_{3}=2.5 \mathrm{~km}^{-1}} \end{gathered}$ | $\begin{gathered} N_{R_{1}=0.6 \mathrm{~km}}-34, \\ N_{R_{2}}=1.2 \mathrm{~km}-1, \\ N_{R_{3}=6.6 \mathrm{~km}}-1 \end{gathered}$ |

## 5. Discussion

Designing of multiple size hexagonal tessellation with minimum overlap difference is a new area in cell planning in telecommunication network design. Our study conjectures an algorithm for efficient masting with least overlap difference for multiple cell range. Application of this formula to MTN River State GSM network solution resulted in an overlap difference of 12.368 km which is a $53.2 \%$ reduction over the original overlap difference of 26.413 km . The formula also uses 36 GSM masts, covering an area of $148.715 \mathrm{~km}^{2}$ compared to the cell engineers original design of 50 masts covering an area of $21.48 \mathrm{~km}^{2}$. This is equivalent to using 1 GSM masts to cover $4.13 \mathrm{~km}^{2}$ in the multiple size hexagonal tessellation model instead of 1GSM masts for $0.43 \mathrm{~km}^{2}$ using the original design. Table 3 shows the results of the computation.

## 6. Conclusion

Our study provide an optimal multiple size hexagonal tessellation design with least overlap difference of 12.368 km and total coverage area of $148.71 \mathrm{~km}^{2}$. The number of GSM masts obtained from the MSHTM is 36 as compared to the original design of 50 GSM masts. This gives a $28 \%$ reduction over the original number of GSM masts. We used geometry of hexagonal tessellation approach to geometric disks covering for multiple cell range to reach optimality and it is the first study that uses multiple size hexagonal tessellation for covering of point sets to arrive at minimum overlap difference and overlap area.

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[^0]:    ${ }^{1}$ Apothem is a line drawn from the center of a regular polygon to an edge and perpendicular to that edge. It is the perpendicular bisector of that edge and also the radius of the inscribed circle to that polygon.

